

Down the rabbit hole of ASMs

①

M - $n \times n$ matrix of entries in an arbit field $\begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{pmatrix}$ Background
System of linear eq's
 $|A| = \det A$ "determines" whether $Ax=b$ has a sol or not.

$$= \sum_{\sigma \in S_n} \text{sgn}(\sigma) a_{1, \sigma(1)} a_{2, \sigma(2)} \dots a_{n, \sigma(n)} \quad (\text{Laplace's formula})$$

Eg $\begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix} = aei - afh - bdi + bfg + cdh - ceg$

where S_n is the set of perms of size n ; i.e. the set of bijective maps $\{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$

Repⁿ: @ Two-line $\begin{pmatrix} 1 & 2 & \dots & n \\ \sigma(1) & \sigma(2) & \dots & \sigma(n) \end{pmatrix} \xrightarrow{\text{One line}} (\sigma(1), \sigma(2), \dots, \sigma(n))$

$\text{sgn}(\sigma) = (-1)^{\text{inv}(\sigma)}$

$\text{inv}(\sigma) = \#$ inversions of $\sigma = \#$ of elem. transp needed to get from σ to $(1, 2, \dots, n)$

$$\hat{\sigma} = |\{i < j : \sigma(i) > \sigma(j)\}|$$

assoc to σ .

② Matrices: $\hat{\sigma}$ is a matrix of 0's & 1's

$$\hat{\sigma}(i, j) = 1 \text{ iff } \sigma(i) = j$$

Many algorithms
→ Laplace exp
→ Cramer's rule
→ Gaussian elim

Eg $n=5, \sigma = 42531, \hat{\sigma} = \begin{pmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix} = \sum_{i < j} \hat{\sigma}_{ij} \hat{\sigma}_{jil}$
 $\text{inv}(\sigma) = 3 + 0 + 1 + 1 = 5$

⇒ Magic square where rows & columns sum to 1.

Then

$$|A| = \sum_{\sigma \in S_n} (-1)^{\text{inv}(\sigma)} \prod_{1 \leq i \leq n} a_{i, \sigma(i)} \quad I, J \subset \{1, \dots, n\} \quad |I| = |J|$$

Dodgson Condensation: Let $A_{i|j}^I$ be the matrix A with i th row, & j th col removed.
 (Lewis Carroll → C. L. Dodgson)
 1866

Desnanot - Jacobi
 (guess 1819) 1835

$$\det A = \frac{\det A_{i|j}^I \det A_{j|i}^n - \det A_{i|i}^I \det A_{j|j}^n}{\det A_{i|i}^n}$$

where $|1| = 1$ if $n=0$
 $\& |a| = a$ if $n=1$ & a

Possible issue if $|A_{1,n}^{1,n}| \neq 0$ (2)

$$\text{Eg) } \begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix} = \frac{\begin{vmatrix} a & b \\ d & e \end{vmatrix} \begin{vmatrix} e & f \\ h & i \end{vmatrix} - \begin{vmatrix} b & c \\ d & f \end{vmatrix} \begin{vmatrix} d & e \\ g & h \end{vmatrix}}{|e|}$$

which is the same.

Nontrivial $\therefore$ Always a polynomial.

λ -Det : ~~Let~~ A - $n \times n$ matrix

$$|A|_{\lambda} := \det_{\lambda} A = \frac{\det_{\lambda} A_1 \det_{\lambda} A_n^{\eta} + \lambda \det_{\lambda} A_n^1 \det_{\lambda} A_1^{\eta}}{\det_{\lambda} A_{1,n}^{1,\eta}}$$

with $|0|_{\lambda} = 1$ & $|a|_{\lambda} = a$.

$$|A|_{-1} = A$$

$$\begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix}_{\lambda} = \frac{\begin{vmatrix} a & b \\ d & e \end{vmatrix}_{\lambda} \begin{vmatrix} e & f \\ h & i \end{vmatrix}_{\lambda} + \lambda \begin{vmatrix} b & c \\ e & f \end{vmatrix}_{\lambda} \begin{vmatrix} d & e \\ g & h \end{vmatrix}_{\lambda}}{|e|_{\lambda}}$$

$$= aei + \lambda(bdj + afh) + \lambda^2(bfg + cdh)$$

$$+ \lambda^3(ceg) + \frac{\lambda(1+\lambda)}{\lambda^2(1+\lambda^{-1})} bde^{-1}f h$$

D. Robbins, H. Rumsey ('86)

$$\begin{pmatrix} 0 & 1 & 0 \\ 1 & -1 & 1 \\ 0 & 1 & 0 \end{pmatrix}$$

$$0+1+0+0+0+0$$

$$|A|_{\lambda} = \sum_{A \in \mathcal{A}_n} \text{inv}(A) (1+\lambda^{-1})^{N(A)} \prod_{i,j} M_{ij}^{A_{ij}}$$

where $\mathcal{A}_n$ is a set of ASMs of size n .

Def An ASM is an $n \times n$ matrix of $0, 1, -1$'s s.t.

a) Rows & columns sum to 1.

b) Non-zero entries alternate in each row.

Eg $\begin{pmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & -1 & 0 & 1 \\ 1 & -1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{pmatrix}$

Facts

- 1) First non-zero entry in every row & col = +1.
- 2) First row & col has one 1.

$\text{inv}(A) = \sum_{\substack{i < j \\ j > k}} A_{ij} A_{kl}$

$N(A) = \# \text{ -1's in } A$

Q) How many ASMs of size n? $|A_n| = ?$

- MRR started to ~~work~~ work on this

- Talked to R. Stanley, who told them that this seq was seen before.

- Andrews solved the ~~CSPs~~ weak Macdonald by counting # of CSPs.

- He defined DPBs & these came up there (later)



* Refined ASMs = Count acc to pos of 1 in $A_{n,k}$

1				
1/2	1	1/2		
2	3/2	3	3/2	2
7	5/2	14	5/2	14

42 2/5 105 7/9 135 9/7 105 5/2 42

429 2/6 1287 8/4 2002 16/16 2002 14/9 1287 6/2 429

Num $\rightarrow (2,1)$ Pascal's Δ le
Denom $\rightarrow (1,2)$ a

Prime factor	n	A _n
	1	1
2	2	2
3	3	7
2, 3, 7	4	42
3, 11, 13	5	429
2, 11, 13	6	7436
2, 13, 17, 19	7	21836
2, 12, 17, 19	8	10850216
2, 5, 17, 19, 23	9	911835460

Round

ASM Conjectures

② $\frac{A_{n,k+1}}{A_{n,k}} = \frac{(n-k)(n+k-1)}{k(2n-k-1)}$ (1986) MRR

① $A_n = \prod_{j=0}^{n-1} \frac{(3j+1)!}{(n+j)!}$

$A_{n,k} = A_{n-1} \frac{(n+k-2)(2n-k-1)}{(n-1)(n-1)}$

NOTE: I could only cover the material up to here.

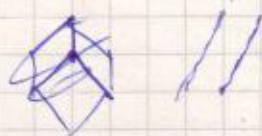
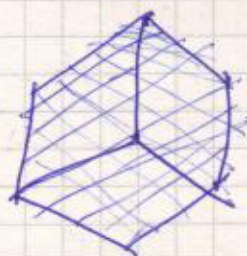
DPPs

Macdonald's cony. (1979)

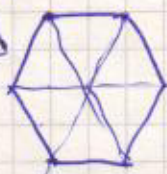
(4)

Andrews proved ~~that~~ # CSPPs :

Lozenge tilings of an $n \times n \times n$ hexagon invariant under 120° rotⁿ in a slab lattice



Handout 1



Eg) $\begin{matrix} 6 & 5 & 5 & 4 & 3 & 3 \\ 6 & 4 & 3 & 3 & 1 & 0 \\ 6 & 4 & 3 & 1 & 1 & 0 \\ 4 & 2 & 2 & 1 & 0 & 0 \\ 3 & 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \end{matrix}$ CSPP of 75

of DPP of order n is a 2D array of positive integers (a_{ij}) defined for $j \geq i$ ~~that has~~ whose rows & columns are weakly decreasing written in the form

$\begin{matrix} a_{11} & a_{12} & a_{13} & \dots & a_{1,\mu_1} \\ & a_{22} & a_{23} & \dots & a_{2,\mu_2} \\ & & \ddots & \ddots & \ddots \\ & & & a_{rr} & a_{r,\mu_r} \end{matrix}$ s.t.

$\begin{matrix} 7 & 7 & 6 & 6 & 3 & 1 \\ & 6 & 5 & 4 & 2 \\ & & 3 & 3 \\ & & & 2 \end{matrix}$

- $\mu_1 \geq \mu_2 \geq \dots \geq \mu_r$
- $a_{ij} \geq a_{i,j+1}$ & $a_{ij} \geq a_{i+1,j}$
- Let $\lambda_i = \mu_i - i + 1 = \#$ elements at row i .
 $a_{ii} > \lambda_i$ & $a_{ii} \leq \lambda_{i-1}$

DPPs of order 3: $\emptyset, \begin{smallmatrix} \times \\ \times \end{smallmatrix}, 2, 3, 31, 32, 33, \begin{smallmatrix} 33 \\ \times \end{smallmatrix}, \begin{smallmatrix} 33 \\ 2 \end{smallmatrix}$

$\mathcal{D}_n :=$ set of DPPs of order n

MRR 1983 / $|\mathcal{D}_n| = A_n$ (Cony.)

But No known bijection

More refinements
AEM DPP

- 1 in first row in pos ~~ok~~
- #1's = m
- Inv# = p

$k-1$ parts equal to n
 m special parts
 p total parts

$a_{ij} \leq j-i$

TSSCPPs (Handout 2)

Count # plane partitions under all possible symmetries

- All 3 reflections
- Rotation
- self-complementary

Handout #2

Let $\mathcal{S}(n)$ = set of TSSCPPs inside a $2n \times 2n \times 2n$ box.

MRR 1986 [Cony] $|\mathcal{S}(n)| = A_n$.

Fundamental region of TSSCPP can be ~~the~~ m_{ij} $1 \leq i \leq n$
 encoded in terms of a Δ le s.t $1 \leq j \leq i$

- $m_{ij} \leq m_{i+1,j}, m_{i,j+1}$
- $m_{ij} \geq j$

eg)

1
 2 3
 4 5 5
 5 5 5 5
 5 5 5 5 5

Magog Δ les

Given ASM, one can ^{also} form the Δ le as follows

1 2 3 4 5 6 7
 1
 1 1 1 1 1
 1 1 1 1 1
 1 1 1 1 1
 1 1 1 1 1
 1 1 1 1 1
 1 1 1 1 1

→ 1 2 3 4 5 6 7
 2 3 4 5 6 7
 1 2 3 4 5 6 7
 1 2 3 4 5 6 7
 1 2 3 4 5 6 7
 1 2 3 4 5 6 7
 1 2 3 4 5 6 7

(g_{ij}) satisfies

$$g_{ij} \leq g_{i-1,j} \leq g_{i,j+1}$$

$$\& g_{ij} < g_{i,j+1}$$

Monotone Δ les = gogis

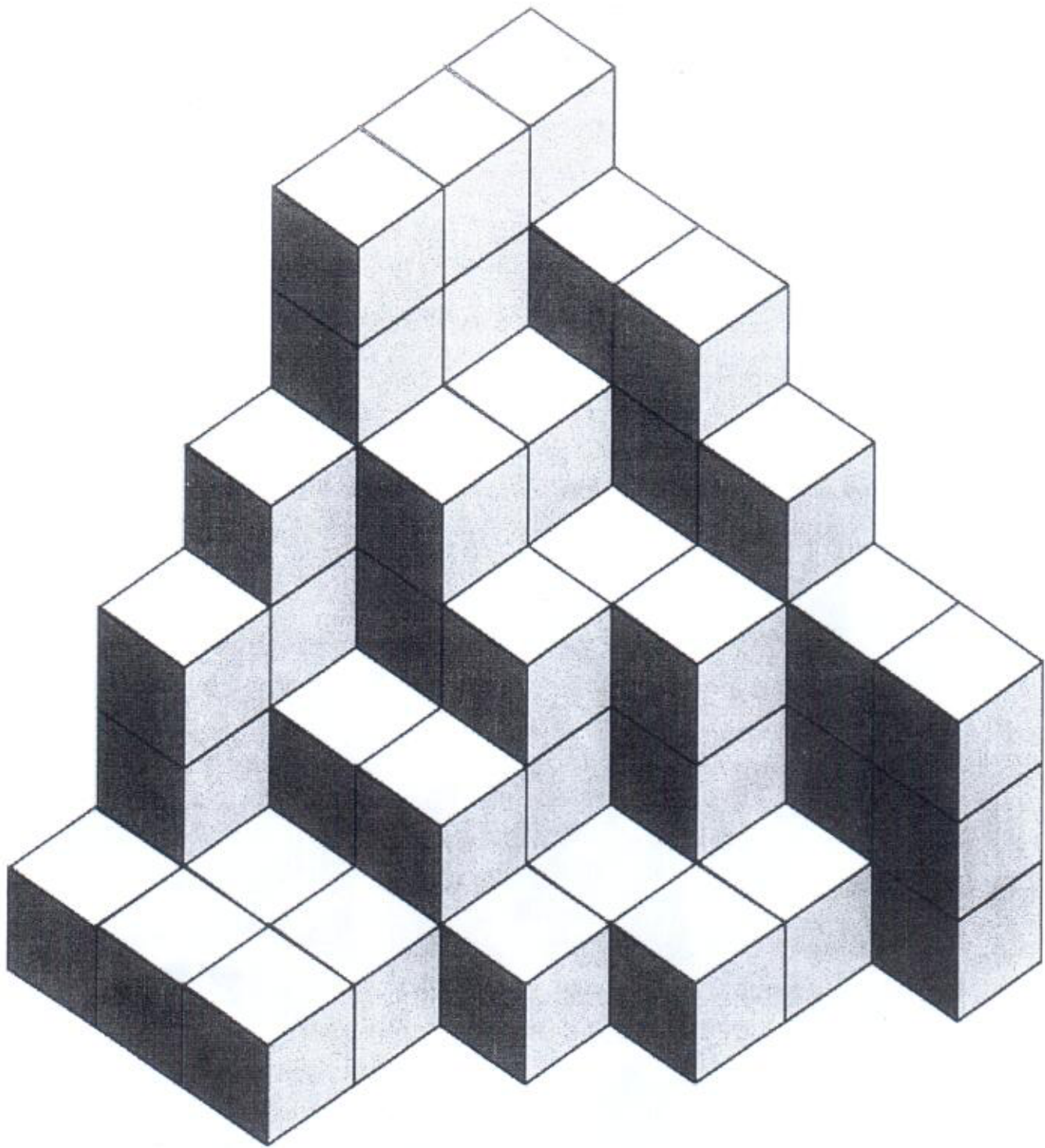
1992 [Zeilberger] - Zeilberger proved gogs = Magogs in a

- page paper with 80 referees

- Kuperberg proved ASM cony in 6 page paper using 6V model (IK det)

- Zeilberger uses it to prove RASM cony.

Refinements?



An example of a CSPP in
a $6 \times 6 \times 6$ box

All 87 TSSCPPs in a
6x6x6 box.

