# **Networking: Hardware and Protocols**

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# Networks

• There are, broadly speaking, two types of networks. Switched and Direct connection.

• Typically, these networks use the PCI bus for interfacing with the computer. Connecting directly to the North Bridge has been proposed for Infiniband.

• Several protocols are in use: TCP/IP, VIA, AM, FM, etc.

#### **Ethernet Network**

• This is the most commonly used network. It is typically used as a switched network, though un-switched configuration can be implemented for a small number of hosts.

• Latency is typically around  $10^2 \ \mu s$ .

• Commonly used implementation operates at 100 Mbps, though Gigabit Ethernet is becoming increasingly popular and 10 Gigabit products are being released.

• Typically this is used with the TCP/IP protocol.

• Using other protocols like VIA or Active Messages can improve performance significantly.

# TCP/IP

• This protocol is used in local area networks.

- Has a layered structure.
- Message is copied across these layers.
- Message is divided into small packets before transmission.
- Path of data is decided at every hop.
- This introduces delays and latency is very poor.

#### **Virtual Interface Architecture**

• Transmission of data is negotiated with the destination. Route of data packets is decided in advance.

- $\circ$  Header has information about all the hops along the way.
- Can be implemented in hardware or emulated in software.
- Is used in Myrinet and Param-net.
- Software emulation for Ethernet is available.

## **Active Messages**

• Direct access to hardware level buffers is provided in user space.

- One sided communications.
- Open source implementation for Linux/Ethernet (GAMMA).
- Can co-exist with TCP/IP.
- $\circ$  Very low latencies can be achieved, e.g.,  $20\mu s$  on Ethernet.
- MPICH has been adapted for GAMMA.

#### **Switched Networks**

- Switched networks offer good bandwidth for multiple concurrent communications. Of course, back-plane of switches plays a crucial role here.
- $\circ$  Ethernet. Bandwidth  $\sim 1Gbps$ , Latency  $\simeq 100 \ \mu s$ .
- $\circ$  Myrinet. Bandwidth  $\sim 2Gbps$ , Latency  $\simeq 5 \ \mu s$ .
- $\circ$  Param-net. Bandwidth  $\sim 2Gbps$ , Latency  $\simeq 10 \ \mu s$ .
- $\circ$  Infiniband. Bandwidth  $\sim 6Gbps$ , Latency  $\simeq 10~\mu s.$

# **Switches**

- Backplane of the switch is very important, unless it is a X-bar switch.
- If a hierarchy of switches is used then switch-switch connections should have a higher bandwidth. Otherwise nodes connected to the same switch have a better connectivity.
- A network is said to have full bisection if any two halves of the network have the same connectivity irrespective of how we divide the network into these parts. Such a network is called Clos network.

# **Switches**

• Switched networks, if backplane is not a bottleneck, offer a good all to all connectivity.

• Number of switches required increases in proportion with (n log(n)).

• Cost of switched networks is often dominated by the cost of switches, though the cost of network interface cards for high performance networks is also very high.

- Rings. (Co-axial cables, SCI/Wulfkit)
  - Each machine connects to two neighbors.
  - $\circ$  Maximum number of hops for a message is n/2.
  - Failure of one node can lead to a breakdown of the network.

- Torus/Mesh. (SCI/Wulfkit)
  - Each machine connects to four neighbors.
  - Maximum number of hops for a message is  $2\sqrt{n}$ .
  - Failure of one node does not disrupt the network.
  - Failure of more than one node disrupts messages.

- 3d Torus. (SCI/Wulfkit)
  - Each machine connects to six neighbors.
  - Maximum number of hops for a message is  $3n^{1/3}$ .
  - Failure of one or two nodes does not disrupt the network.
  - Failure of more than two nodes disrupts messages.

- Hyper-Cube.
  - Machines are connected along a k dimensional cube,  $n = 2^k$ .
  - $\circ$  There are at most k hops between nodes.
  - $\circ$  A total of k\*n cables are needed.

#### • Fully connected.

 $\circ$  All machines are connected to each other.

 $\circ$  A total of  $n^2/2$  cables are needed.

 $\circ$  For n = 128, we require 8192 cables.

- Topologies mentioned here can also be realized using switched networks.
- Routing in multiply connected networks is a complex problem. • As long as the number of hops is small (~ 10), switch-less technologies like SCI/Wulfkit offer good performance. These have a latency of  $5\mu$ s and a point to point bandwidth of more than 2Gbps.
- $\circ$  Cost of these networks grows in proportion with the number of nodes.

#### Networks

• There are other solutions like Quadrics, Cray interconnect, etc. that have better performance than the solutions discussed here.

• Most networking solutions used here interface through the PCI bus. This will restrict bandwidth to less than 8 Gbps in near future.

• Infiniband can potentially interface through the memory bus directly. This will reduce latency by a significant amount and also remove the bottleneck on bandwidth.

#### Networks

• In last 30 years CPU speed has improved by a factor of 3,000.

In the same period, memory latency has improved by a factor of 30.
This lead to the development of cache in order to hide the widening gap in performance of memory and CPUs.

In the same period, latency has improved by a factor of 10 for
Ethernet. For networks in general, it has improved by a factor of 50.
How will technology be used to address this gap further?