

Systems Biology Across Scales: A Personal View

XII. Intra-cellular Systems II: Metabolism and Modularity

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Intra-cellular biochemical networks

□ **Metabolic networks**

Nodes: metabolites (substrates & products of metabolism)

Links: chemical reactions (directed)

□ **Genetic regulatory networks**

Nodes: Genes & Proteins

Links: regulatory interactions (directed)

□ **Protein-Protein interaction network**

Nodes: Proteins

Links: physical binding and formation of protein complex (undirected)

□ **Signaling network**

Nodes: Signaling molecules e.g., kinase, cyclicAMP, Ca

Links: chemical reactions (directed)

Metabolism

Chemical process through which cells break down nutrients to generate energy and/or into usable building blocks (*catabolic metabolism*) and then reassemble them using energy to form biological molecules necessary for the cell (*anabolic metabolism*).

Uses sequence of chemical reactions (pathways) to convert *substrates* (initial inputs) successively into useful *products*.
Reactions are aided by *enzymes*.

Metabolic network: The set of all reactions in all pathways

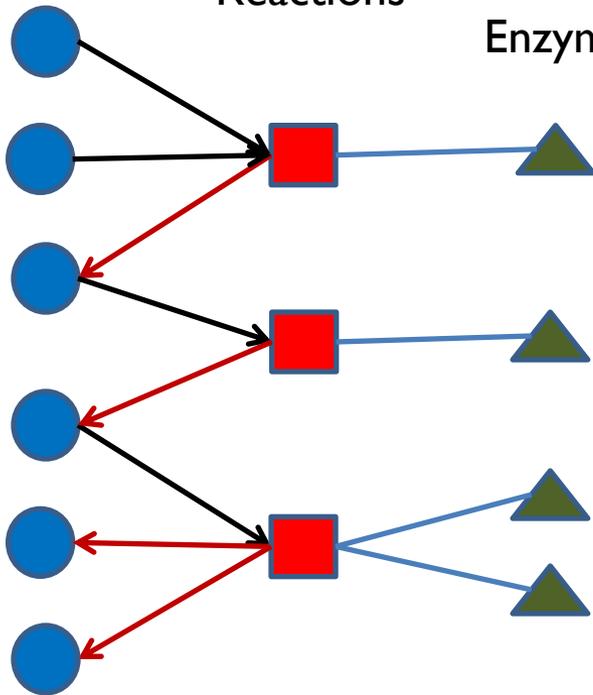
Representing Metabolic Networks

Tripartite directed network

Metabolites

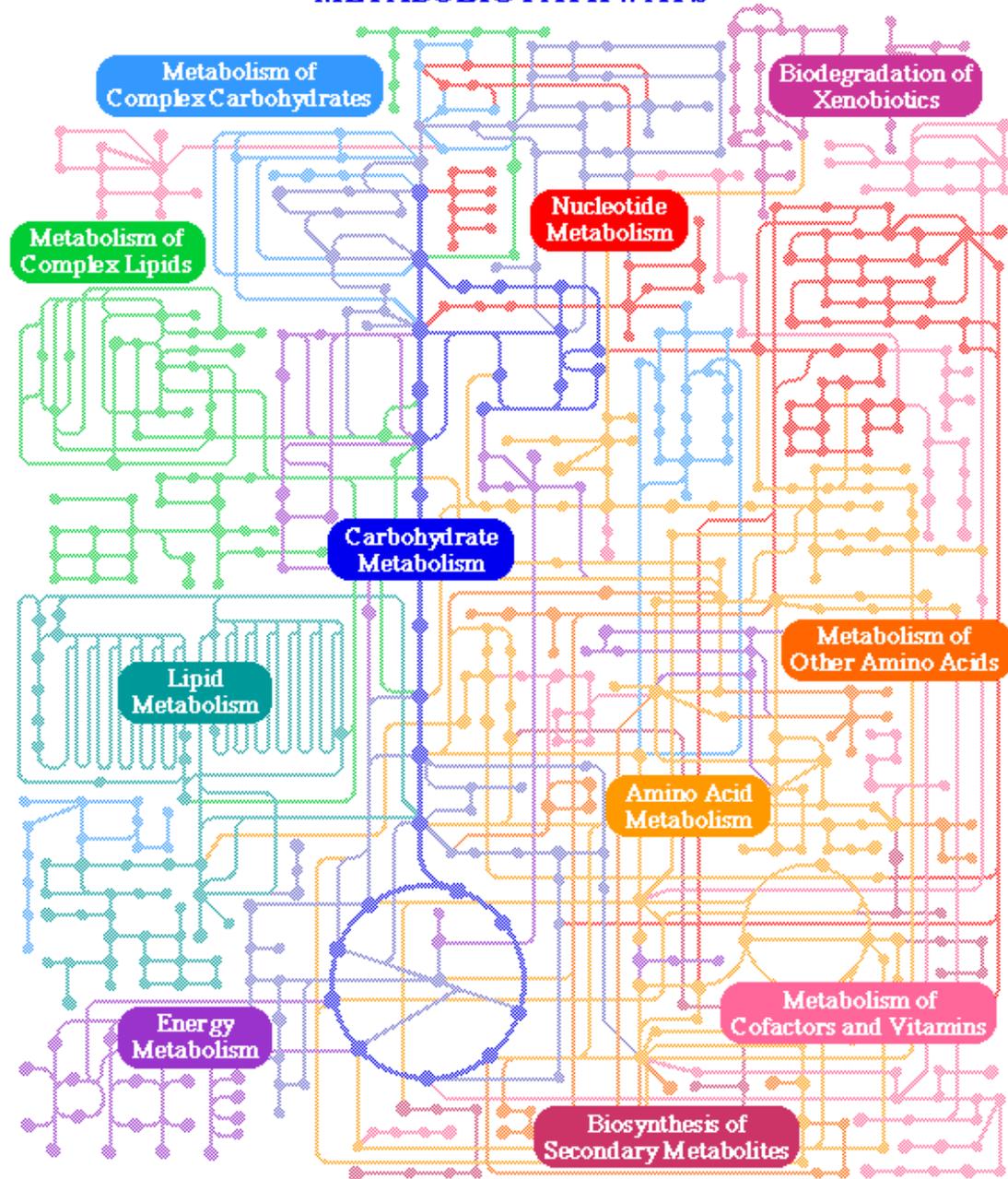
Reactions

Enzymes



Projection to only metabolites

METABOLIC PATHWAYS



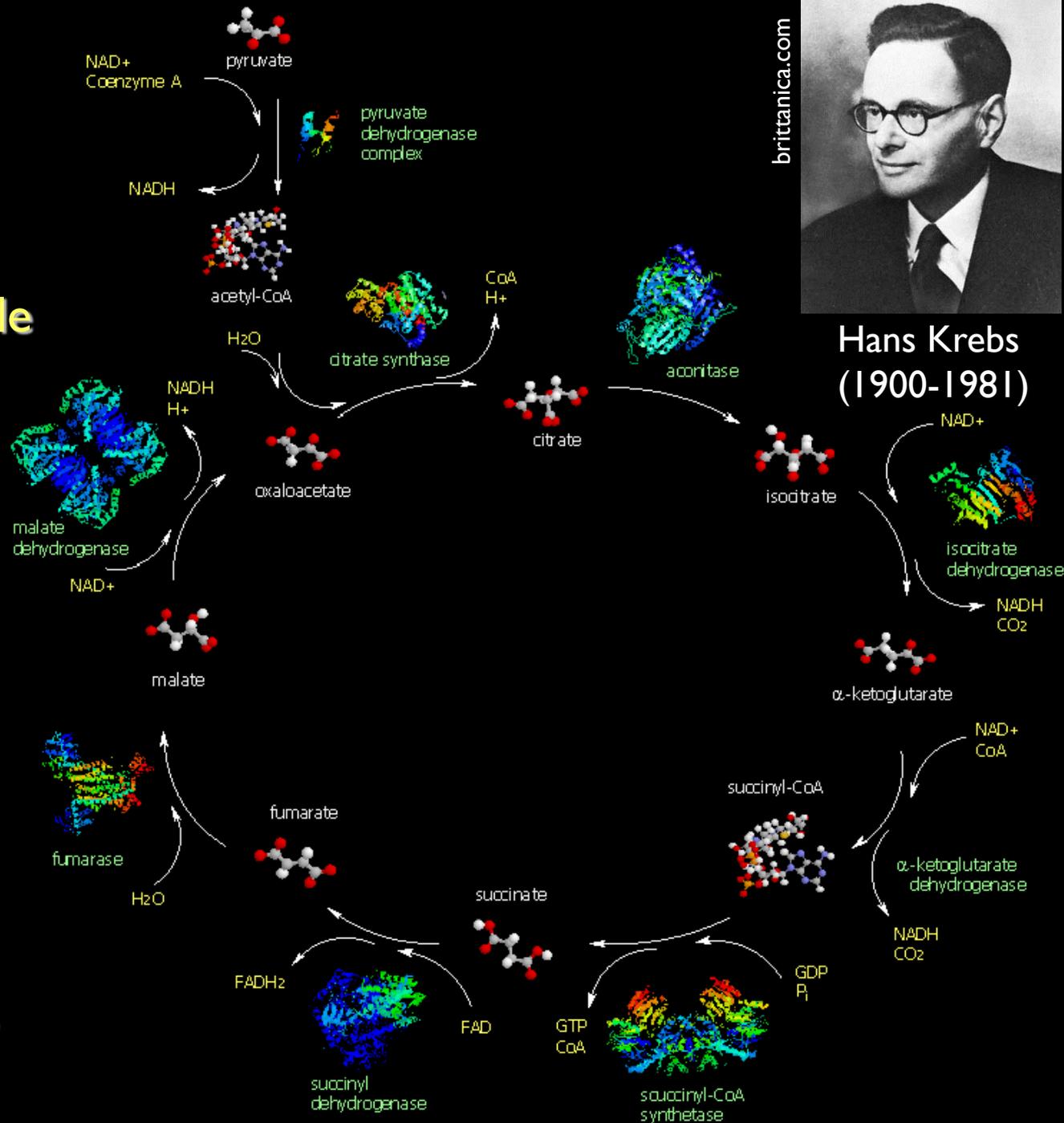
KEGG

Citric acid cycle

Also known as Tricarboxylic Acid (TCA) or Krebs cycle

is a series of enzyme-catalysed chemical reactions lying at the heart of aerobic metabolism.

Involved in the breakdown of all 3 major food groups: carbohydrates, lipids and proteins.

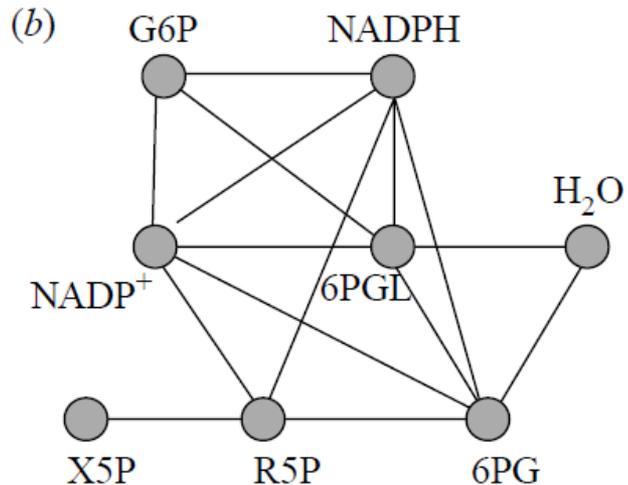
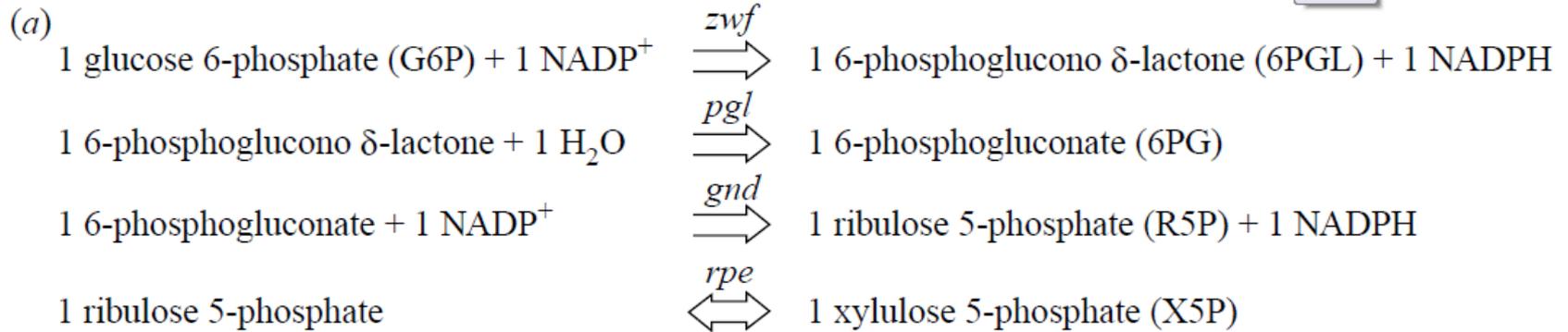


britannica.com

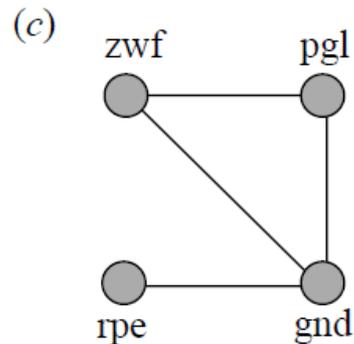


Hans Krebs (1900-1981)

Graphical representation of metabolic networks



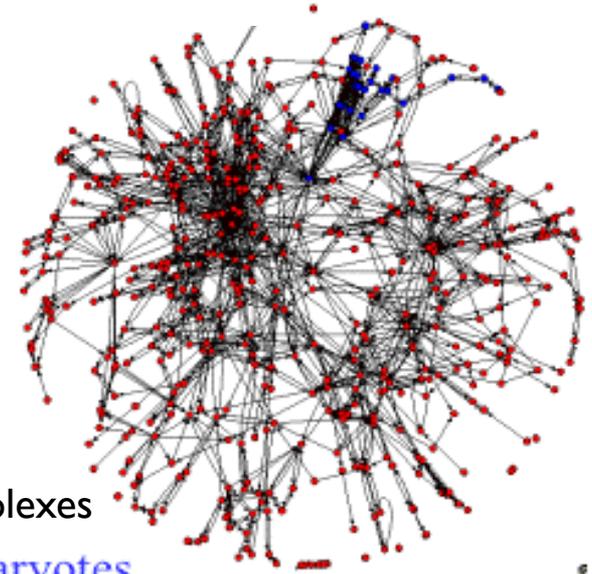
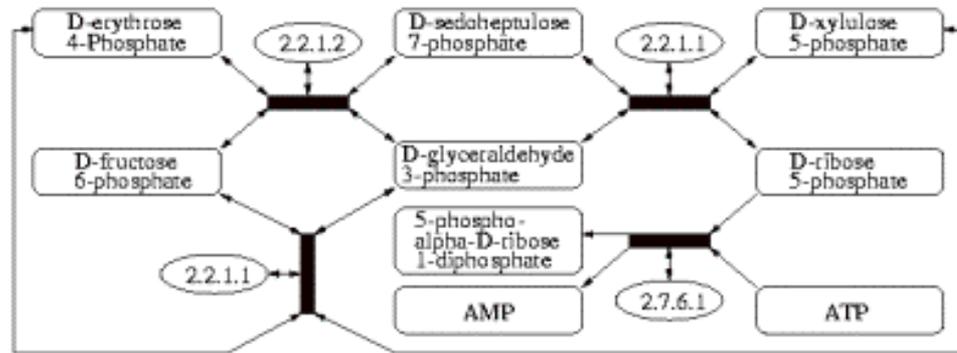
Substrate graph



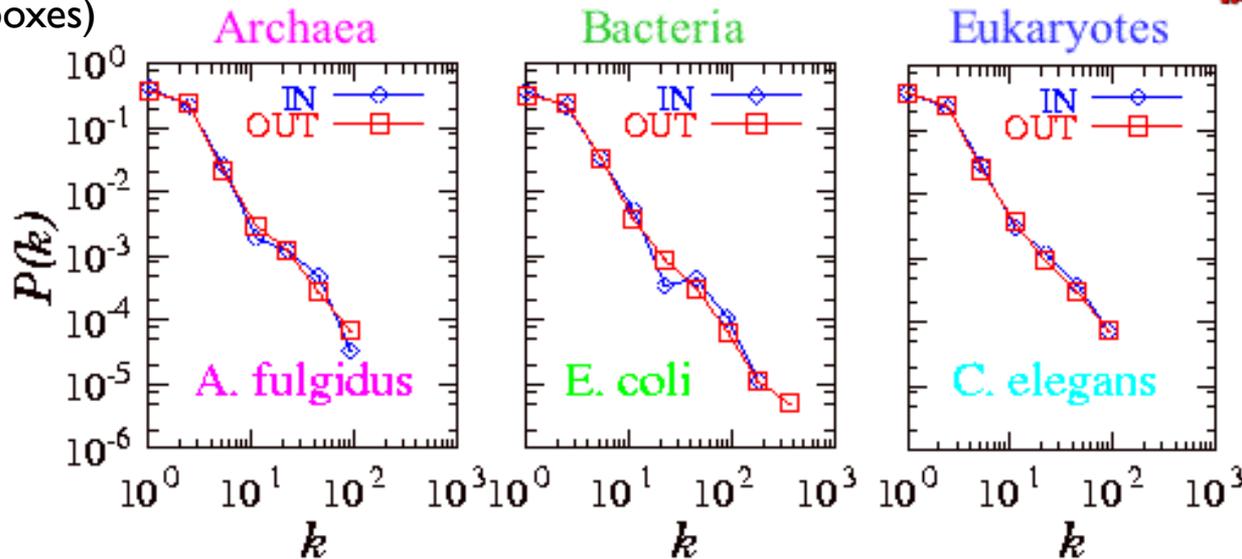
Reaction graph

Scale-free nature of degree distribution of metabolic networks

A portion of the WIT database for *E. coli*.



Nodes are substrates & products, linked by enzyme-substrate complexes (black boxes)



Organisms from all three domains of life are **scale-free** networks!

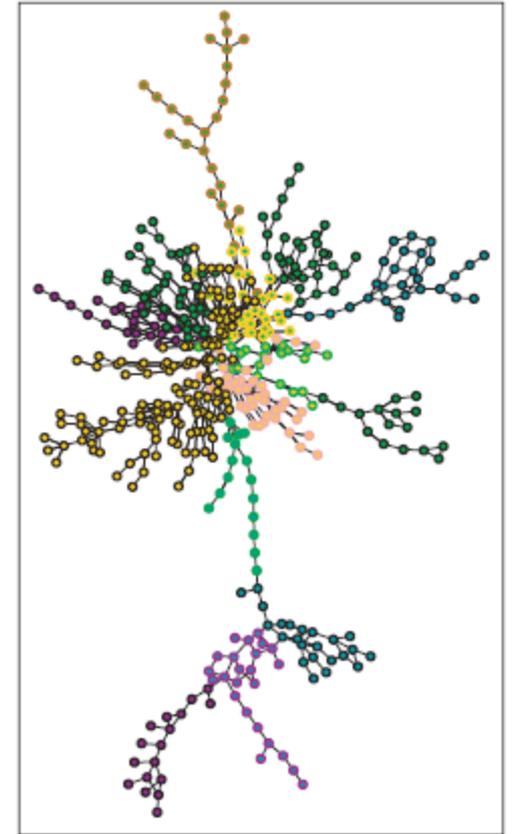
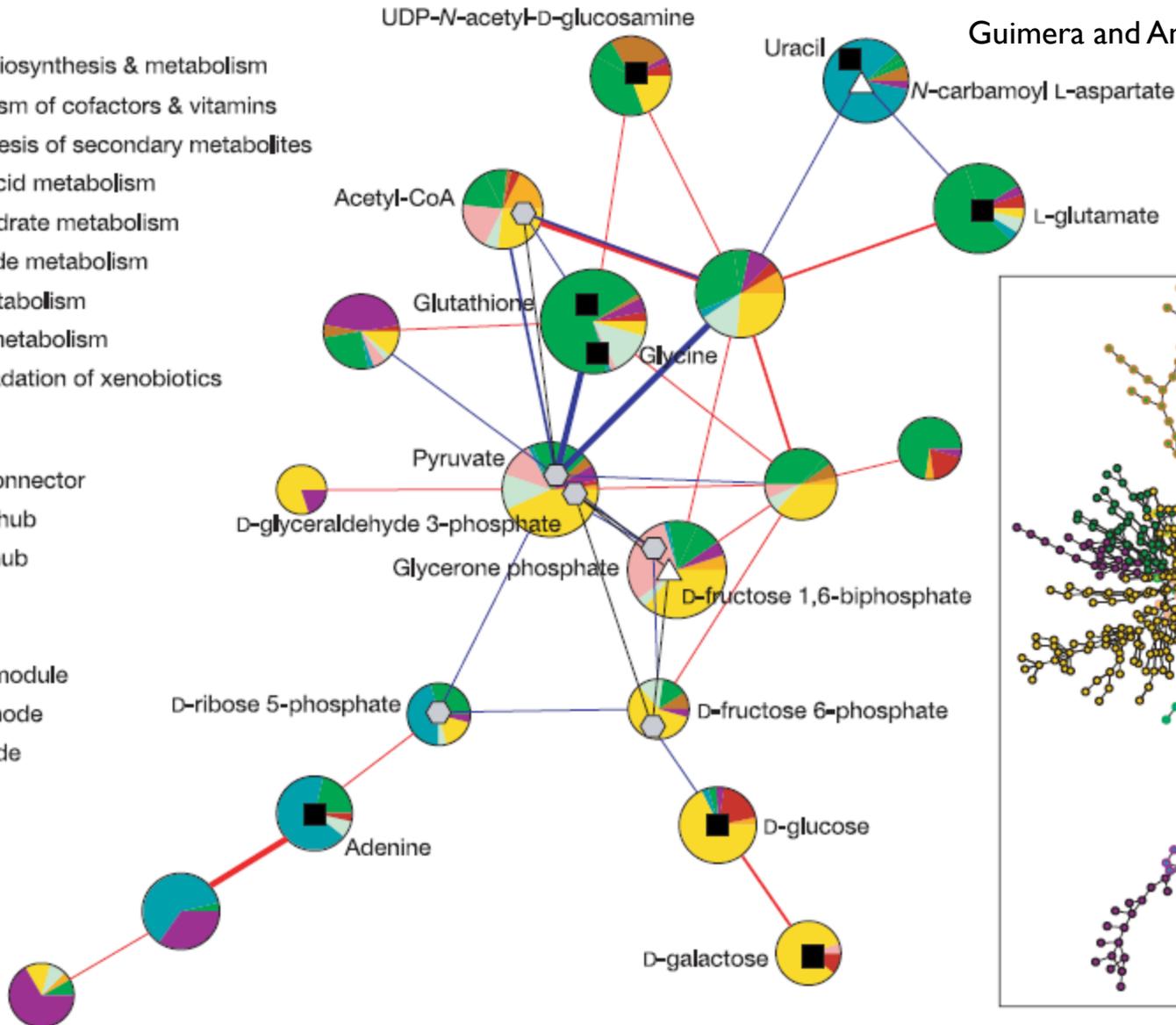
Modular nature of metabolic networks

Guimera and Amaral, *Nature* (2005)

-  Glycan biosynthesis & metabolism
-  Metabolism of cofactors & vitamins
-  Biosynthesis of secondary metabolites
-  Amino-acid metabolism
-  Carbohydrate metabolism
-  Nucleotide metabolism
-  Lipid metabolism
-  Energy metabolism
-  Biodegradation of xenobiotics

-  Non-hub connector
-  Connector hub
-  Provincial hub

-  Module-module
-  Module-node
-  Node-node



Each circle represents a module and is colored according to the KEGG pathway classification of the metabolites it contains.

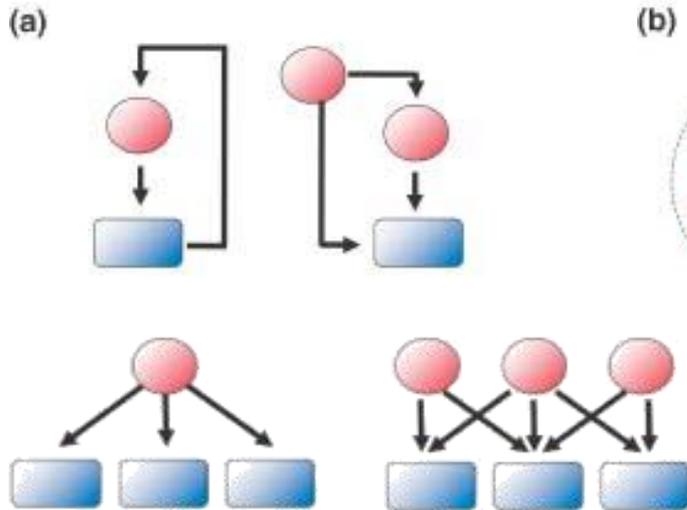
Metabolic network of *E. coli* (N=473, L =674).

Modular Networks: dense connections *within* certain sub-networks (**modules**) & relatively few connections *between* modules

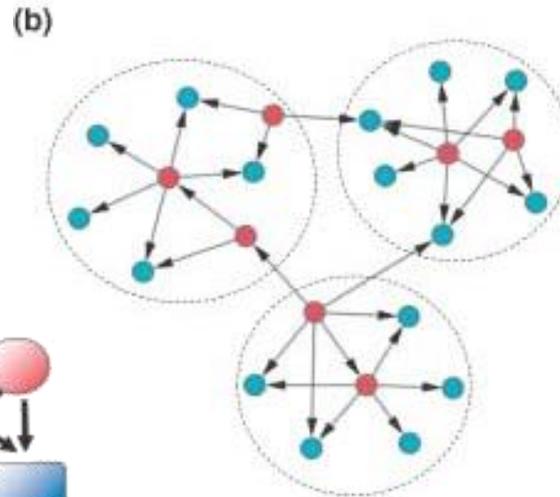
Modules: A *mesoscopic* organizational principle of networks

Going beyond *motifs* but more detailed than *global* description (L, C etc.)

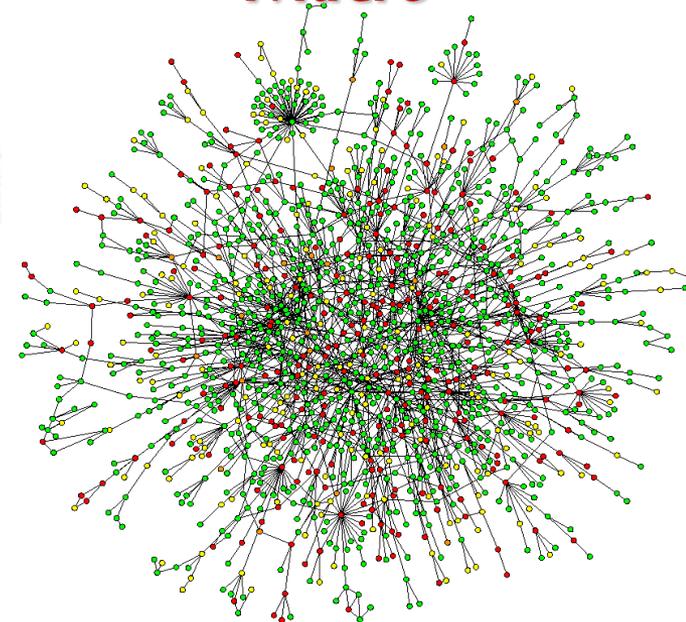
Micro



Meso



Macro



Modular Biology (Hartwell et al, Nature 1999)

Functional modules as a critical level of biological organization

Modules in biological networks are often associated with specific functions

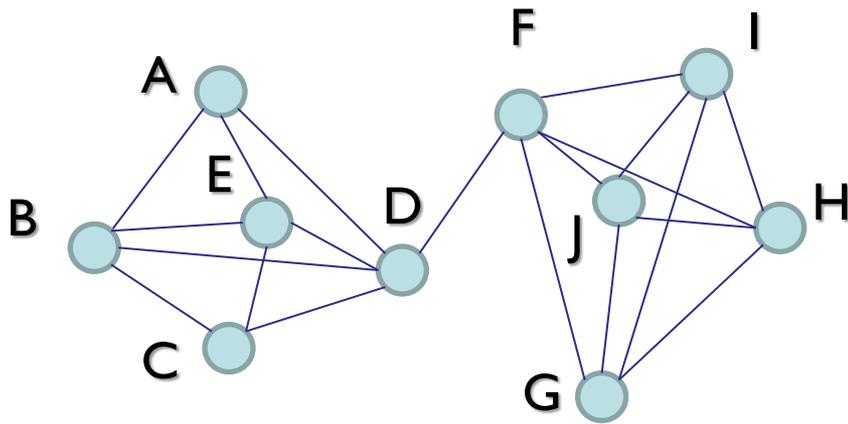
Problem:

Given a network,
how do we find the modules (communities)
into which it can be divided ?

Community Detection in Networks

Also referred to as Graph Partitioning or Module Determination

How to divide the nodes of a network into several groups such that nodes in each group are densely or strongly inter-connected



E.g., it is clear that node clusters I: {A,B,C,D,E} and II: {F,G,H,I,J} constitute two separate groups that are highly intra-connected but has only a single link connecting the two groups

The corresponding adjacency matrix will have an almost block-diagonal form – the two blocks corresponding to node clusters I & II

However for large networks the modular character may not be visually apparent – and adjacency matrices need to be partitioned

Graph partitioning

A classic problem in computer science from 1960s

How to divide the nodes of a network into a given number of non-overlapping groups of given sizes such that the number of edges between groups is minimized ?

A generalization of this problem,

How to divide the nodes into several groups such that most links are within groups and few links are between groups

referred to as

Community detection

How we define “most” and “few” can vary from one algorithm to another

Spectral partitioning

Fiedler 1973

Consider a network of N nodes and L links

Aim: to divide the N nodes into 2 groups (Groups A and B, say) to reduce the **cut size** (number of links between the two groups)

$R = (1/2) \sum_{ij} A_{ij}$ such that **i and j belong to different groups**

Partitioning into more than 2 groups can be done by **repeated bisection**

For each node, a label $s = \{-1, +1\}$ is defined

$s_i = +1$ if node i belongs to group A, $= -1$ if i belongs to group B

Thus

$(1/2) (1 - s_i s_j) = 1$ if i & j are in different groups,
 $= 0$ if i & j are in same group

$$\Rightarrow R = (1/4) \sum_{ij} A_{ij} (1 - s_i s_j) = (1/4) \sum_{ij} (k_i \delta_{ij} - A_{ij}) s_i s_j$$

$$\Rightarrow R = (1/4) \sum_{ij} L_{ij} s_i s_j \text{ where } \mathbf{L} = \mathbf{D} - \mathbf{A} \text{ is the Laplacian matrix}$$

In matrix notation $R = (1/4) \mathbf{s}^T \mathbf{L} \mathbf{s}$ where $\mathbf{s} = \{s_1 s_2 \dots s_N\}$

Goal of Partitioning: To find \mathbf{s} that minimizes R given L

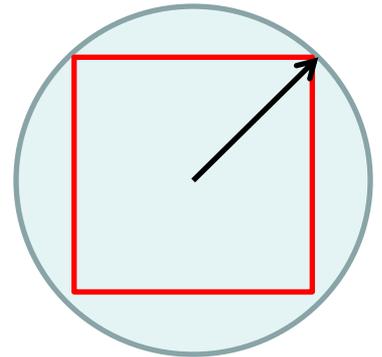
Partitioning as minimization

Fiedler 1973

If s_i were allowed to take any possible value, then differentiation gives the optimum

But s_i are restricted to $\{-1, +1\} \Rightarrow$ difficult problem

\mathbf{s} can be seen as a vector that points to any one of the 2^N vertices of N -dimensional hypercube



Possible approximate solution:

Allow s_i to take any value subject to the constraints that

- (i) $\sum_i s_i^2 = N \Rightarrow \mathbf{s}$ is a vector in N -dimensional unit hypersphere
- (ii) $\sum_i s_i = N_A - N_B$ where N_A, N_B are the sizes of the two groups

In matrix notation $\mathbf{1}^T \mathbf{s} = N_A - N_B$

The minimization problem can now be solved as

$$\frac{\partial}{\partial s_i} [\sum_{jk} L_{jk} s_j s_k + \lambda (N - \sum_j s_j^2) + 2\mu ([N_A - N_B] - \sum_j s_j)] = 0$$

where λ, μ are Lagrange multipliers for enforcing the constraints

$$\Rightarrow \sum_j L_{ij} s_j = \lambda s_i + \mu \quad \text{In matrix notation, } \mathbf{L} \mathbf{s} = \lambda \mathbf{s} + \mu \mathbf{1}$$

Partitioning using the Laplacian spectrum

Multiplying $\mathbf{L} \mathbf{s} = \lambda \mathbf{s} + \mu \mathbf{1}$ by $\mathbf{1}^T$ on the left

we get $\lambda [N_A - N_B] + \mu N = 0 \Rightarrow \mu/\lambda = -[N_A - N_B]/N$

Using $\mathbf{1}^T \mathbf{s} = N_A - N_B$ and $\mathbf{1}^T \mathbf{L} = 0$ ($\mathbf{1}$ is eigenvector of L with eigenvalue 0)

Defining a new vector $\mathbf{x} = \mathbf{s} + (\mu/\lambda) \mathbf{1} = \mathbf{s} - \mathbf{1}[N_A - N_B]/N$

$\Rightarrow \mathbf{L} \mathbf{x} = \mathbf{L} (\mathbf{s} + (\mu/\lambda) \mathbf{1}) = \mathbf{L} \mathbf{s} = \lambda \mathbf{s} + \mu \mathbf{1} = \lambda \mathbf{x}$

Thus \mathbf{x} is an eigenvector of the Laplacian

But which eigenvector ?

The one that gives the smallest value of cut size R

We can't choose $\mathbf{1} = \{1, 1, \dots, 1\}$ as it is orthogonal to \mathbf{x} because $\mathbf{1}^T \mathbf{x} = 0$

Note that cut size is proportional to the eigenvalue λ

as $R = (1/4) \mathbf{s}^T \mathbf{L} \mathbf{s} = (1/4) \mathbf{x}^T \mathbf{L} \mathbf{x} = (1/4) \lambda \mathbf{x}^T \mathbf{x} = \lambda [N_A N_B]/N$

Thus we have to choose the eigenvector corresponding to the lowest non-zero eigenvalue (smallest eigenvalue of L is 0 with eigenvector $\mathbf{1}$)

Finally, optimal partition \mathbf{s} is obtained from $\mathbf{s} = \mathbf{x} + \mathbf{1}[N_A - N_B]/N$

From Relaxation Approxn to the Network

For the actual network, the optimal partition \mathbf{s} is subject to the additional constraint that (i) $s_i = +1$ or -1 , and (ii) exactly N_A of the components are $+1$ and N_B are -1

Thus, we need to choose \mathbf{s} as close as possible to ideal value subject to the constraints \Rightarrow maximize the vector length, i.e.,

$$\mathbf{s}^T \mathbf{s} = \mathbf{s}^T (\mathbf{x} + \mathbf{1} [N_A - N_B]/N) = \sum_i s_i (x_i + [N_A - N_B]/N)$$

by assigning $s_i = +1$ for the nodes corresponding to the N_A largest (most positive) values of \mathbf{x} , i.e., the components of the eigenvector of the lowest non-zero eigenvalue of \mathbf{L} , and, $s_i = -1$ to the remaining N_B nodes

Note: If $N_A \neq N_B$, we can either choose (i) N_A elements to be $+1$ (N_B elements -1) or (ii) N_A elements to be -1 (N_B elements $+1$)
The one having lower cut size is the optimal partition

Community detection

How to quantify the degree of modularity for a given partitioning of a network into communities ?

Is there a distinction between links within a module and that between a module and the rest ?

One suggested measure:

$$Q \equiv \frac{1}{2L} \sum_{ij} \left[A_{ij} - \frac{k_i k_j}{2L} \right] \delta_{c_i c_j} \quad (\text{Newman, EPJB, 2004})$$

= 1 if nodes are in same community

probability of an edge betn 2 nodes proportional to their degrees

A: Adjacency matrix

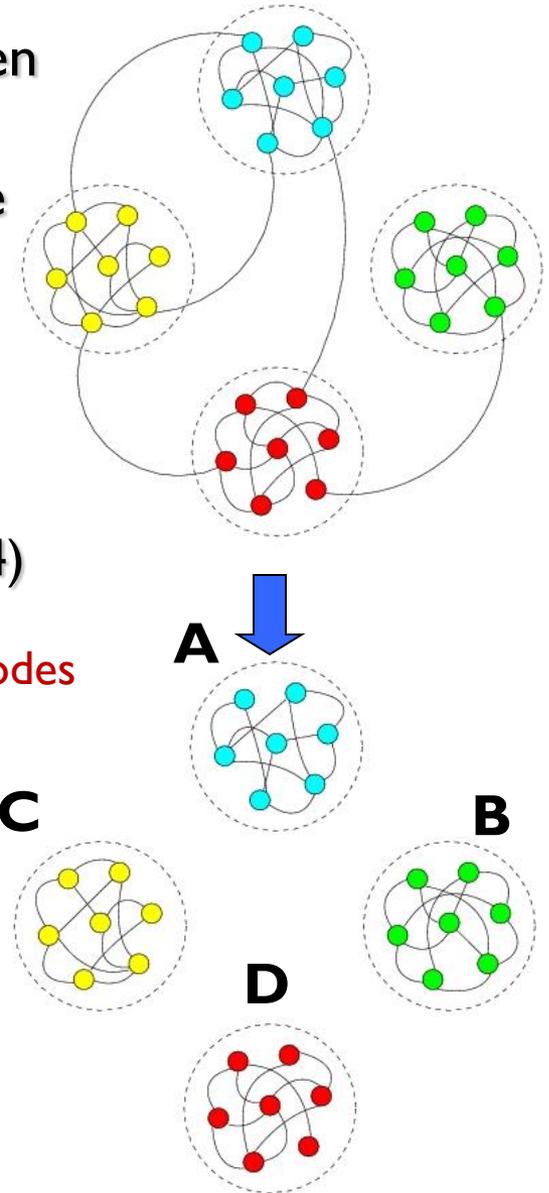
L : Total number of links

k_i : degree of i -th node

c_i : label of module to which i -th node belongs

For a random network, $Q = 0$

i.e., the connection density within a module is no different from that anywhere else in the network



Community detection

For directed & weighted networks:

$$Q^W \equiv \frac{1}{L^W} \sum_{ij} \left[W_{ij} - \frac{s_i^{\text{in}} s_j^{\text{out}}}{L^W} \right] \delta_{c_i c_j} \quad (L^W = \sum_{ij} W_{ij})$$

W: Weight matrix

s_i : strength of i -th node

Modules determined through a generalization of the spectral method (Leicht & Newman, 2008)

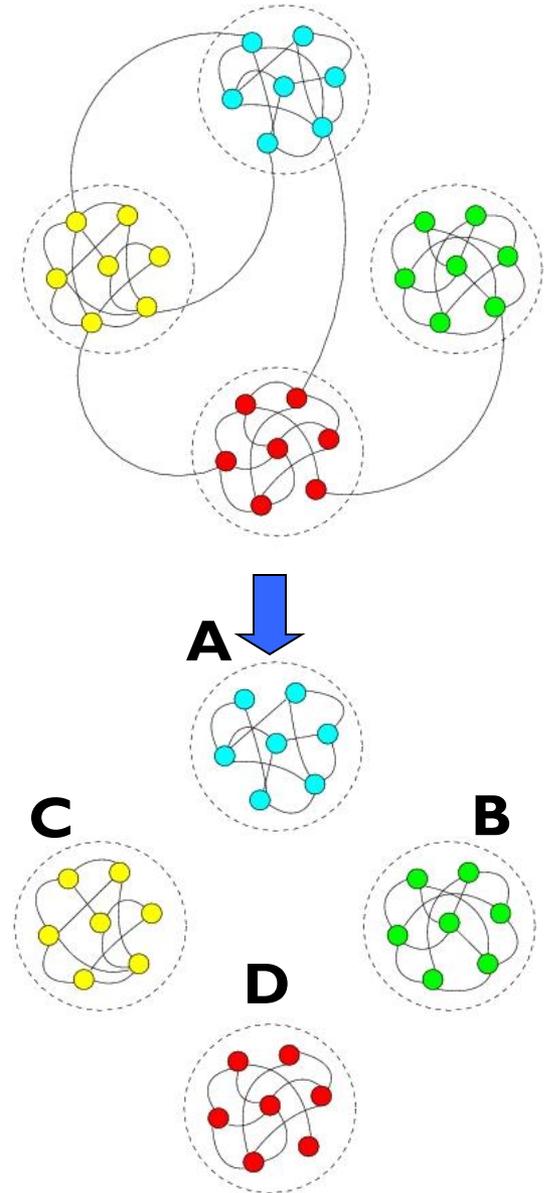
Calculate eigenvector corresponding to largest +ve eigenvalue of symmetrized modularity matrix

$\mathbf{B} + \mathbf{B}^T$ where

$$B_{ij} = W_{ij} - [s_i^{\text{in}} s_j^{\text{out}} / L^W]$$

and then assign communities based on the signs of the elements of the eigenvector.

Simplest generalization of the method to more than 2 communities is to use repeated bisection

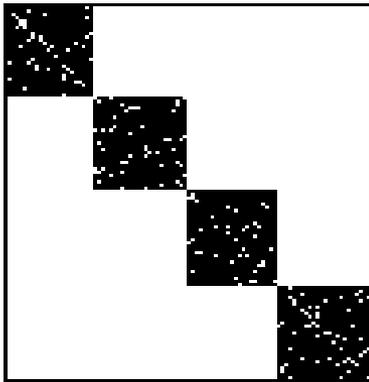


A simple model of modular networks

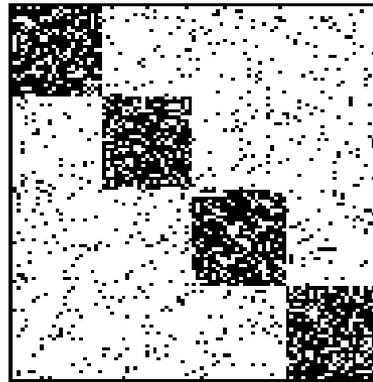
Model parameter r :

Ratio of inter- to intra-modular connection density

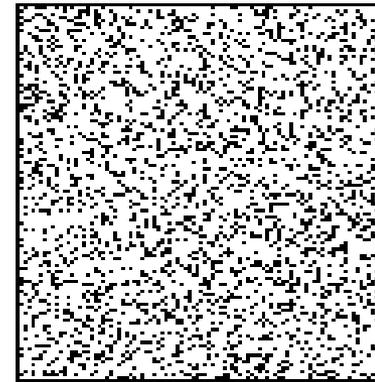
(a) $r = 0$



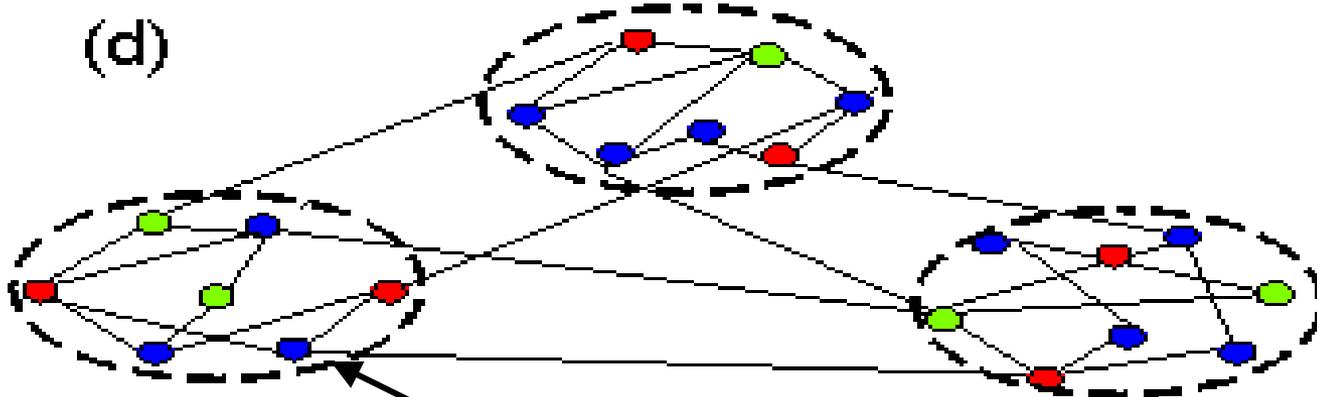
(b) $r = 0.1$



(c) $r = 1$



(d)



Module \equiv random network

Comparison with Watts-Strogatz model

Structural measures used:

Communication efficiency

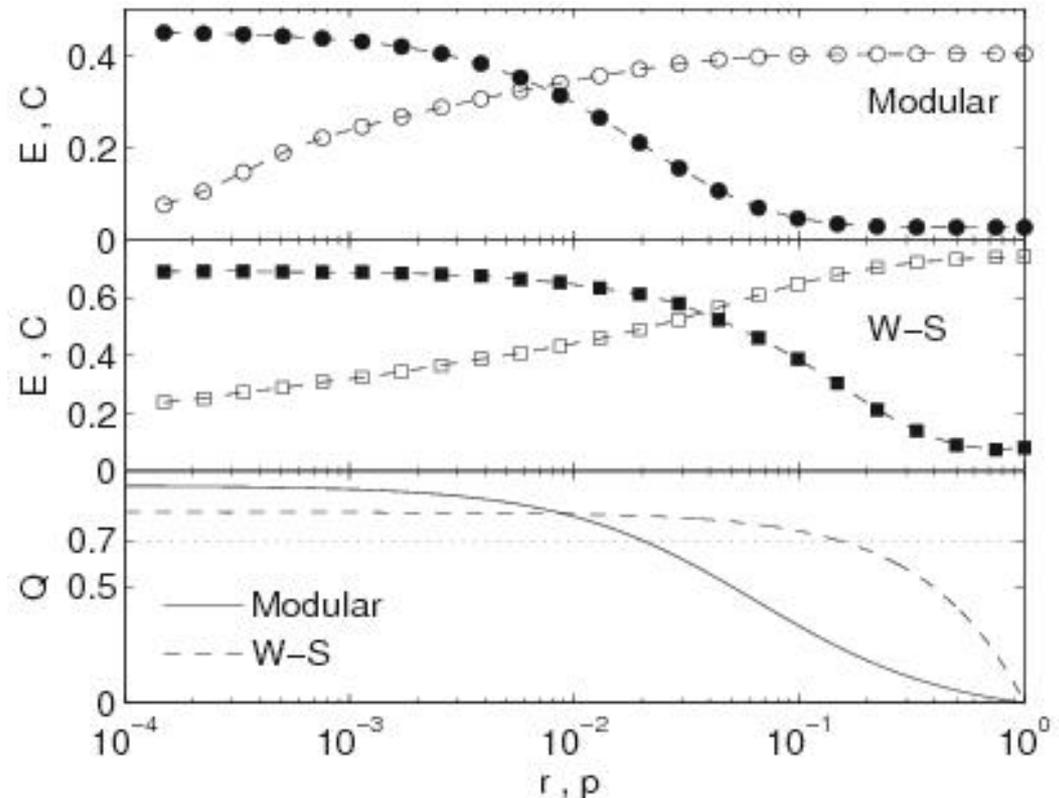
$$E = [\text{avg path length, } \ell]^{-1} = 2 / N(N-1) \sum_{i>j} d_{ij}$$

Clustering coefficient

$$C = \text{fraction of observed to potential triads} \\ = (1 / N) \sum_i 2n_i / k_i (k_i - 1)$$

WS and Modular networks behave similarly as function of p or r (Also for between-ness centrality, edge clustering, etc)

In fact, for same N and $\langle k \rangle$, we can find p and r such that the WS and Modular networks have the same “modularity” Q



How can you tell them apart ?

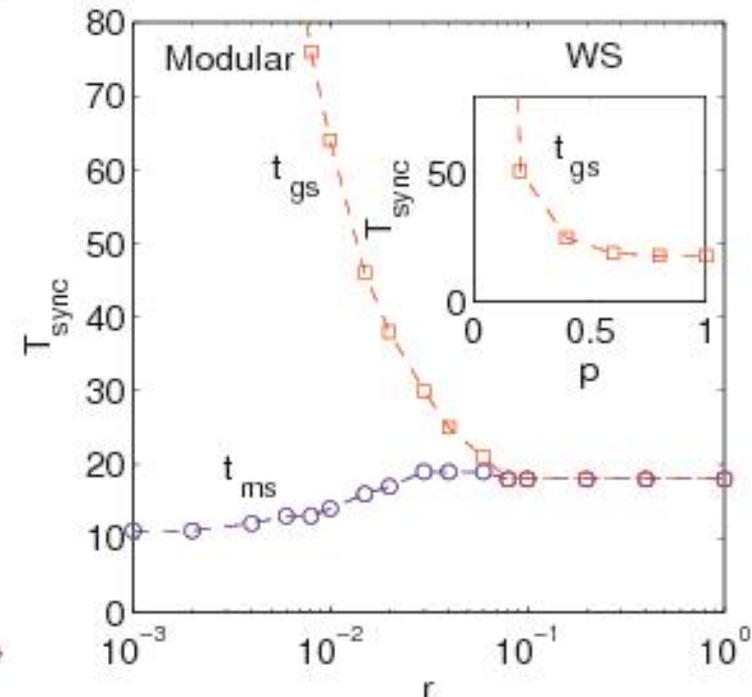
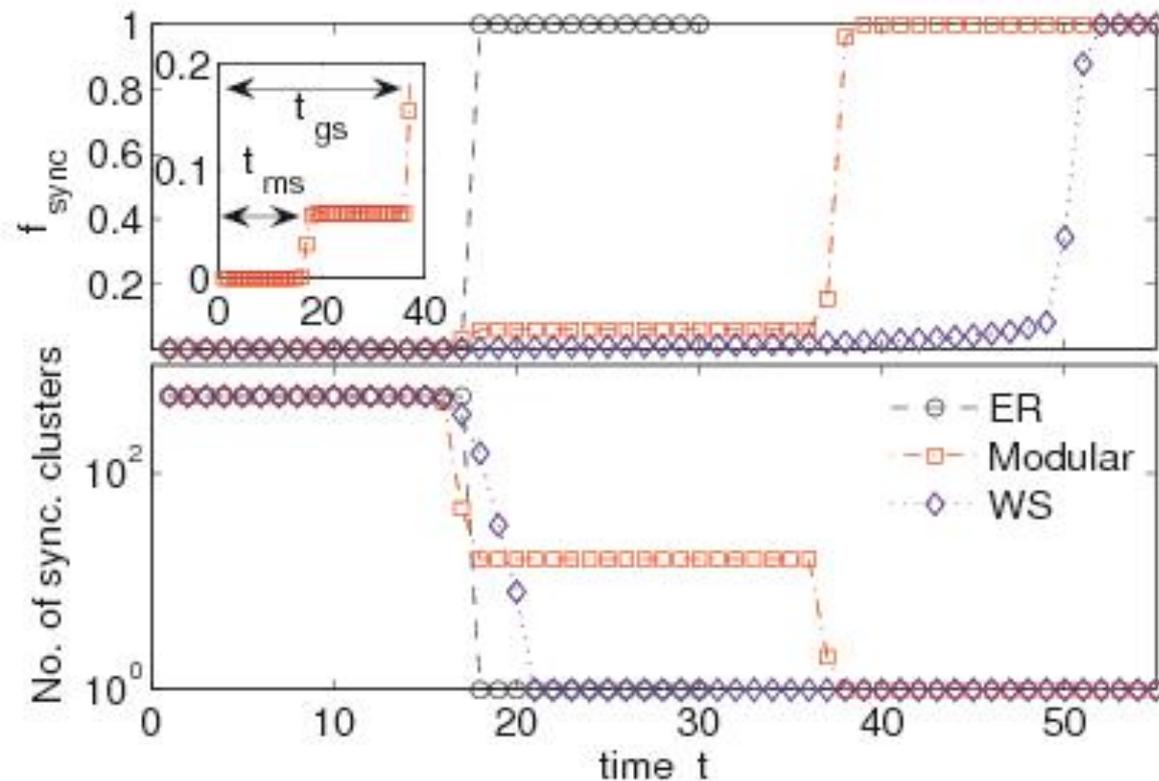
Dynamics on Modular networks different from that on Watts-Strogatz small-world networks

Consider synchronization on modular networks

e.g., phase oscillators: $d\theta_i/dt = \omega + (1/k_i) \sum K_{ij} \sin(\theta_j - \theta_i)$

Network topology

2 distinct time scales in Modular networks: t_{modular} & t_{global}



Existence of distinct time-scales in Modular networks

Pan and Sinha, EPL (2009)

Consider linearized dynamics around synchronized state

$$d\theta_i / dt = - (\kappa / k_i) \sum_j L_{ij} \theta_j, \quad (i = 1, \dots, N)$$

L: Laplacian

Focus on the normal modes:

$$\phi_i(t) = \sum_j B_{ij} \theta_j = \phi_i(0) \exp(-\lambda_i t), \quad (i = 1, \dots, N)$$

κ : coupling strength of oscillators

B: matrix of eigenvectors
 λ_i : eigenvalues

} of $L' = D^{-1} L$,
D: diagonal matrix s.t. $D_{ii} = k_i$

$L' \rightarrow L = D^{1/2} L' D^{-1/2}$ is symmetric, normalized Laplacian $\Rightarrow \lambda_i$ real

Differences in time-scales of modes \Rightarrow gap in spectrum of L

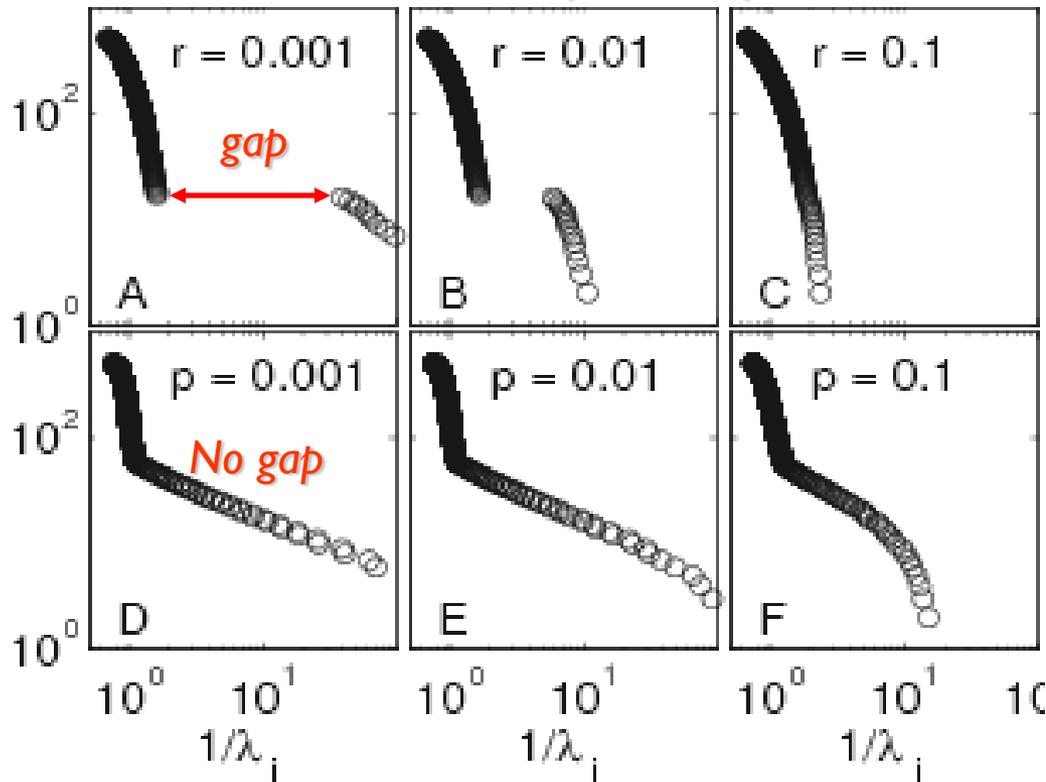
Mode for smallest λ_i : associated with global synchronization

Other modes: synchronization within different groups of oscillators

Eigenvalue spectra of the Laplacian

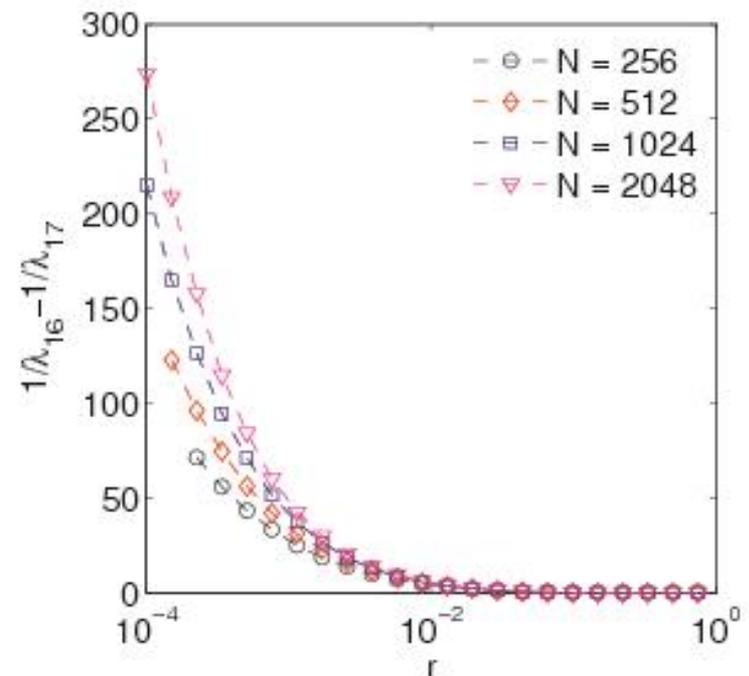
Shows the existence of spectral gap \Rightarrow distinct time scales

Modular network Laplacian spectra



WS network Laplacian spectra

Spectral gap in modular networks diverges with decreasing r



Existence of distinct time-scales in Modular networks

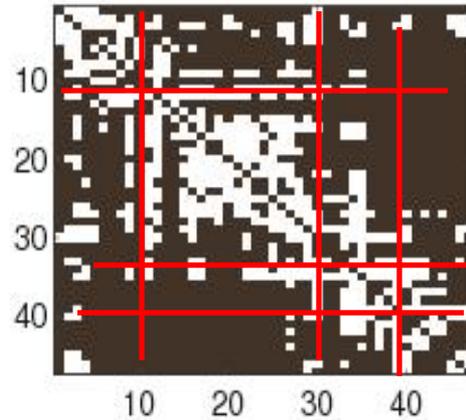
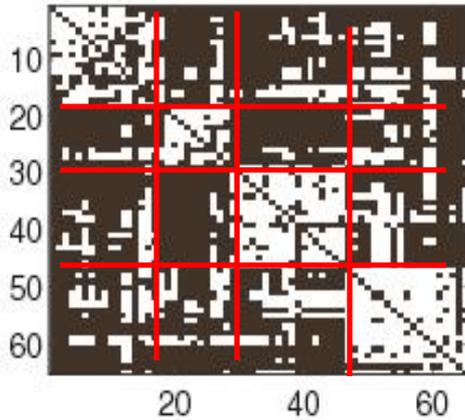
No such distinction in Watts-Strogatz small-world networks

How about “real” SW networks ?

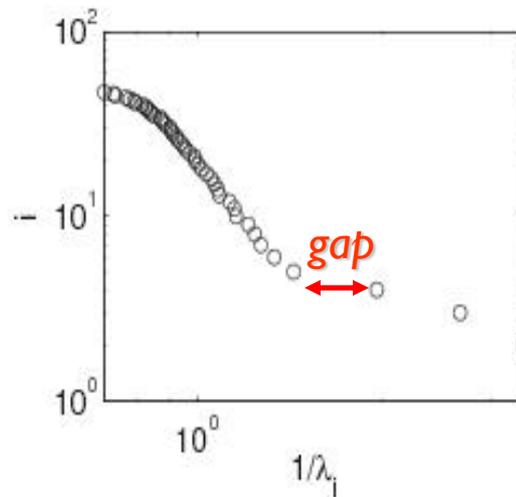
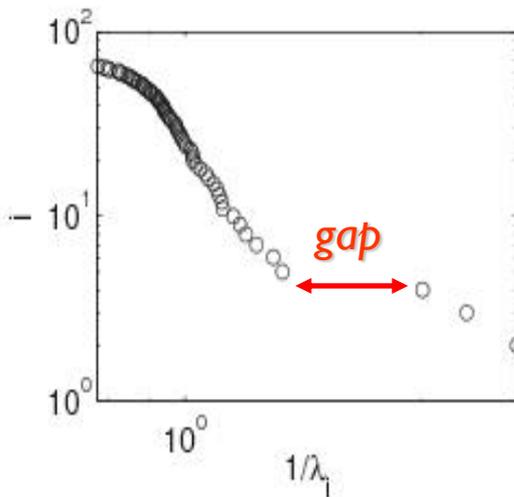
Pan and Sinha, EPL (2009)

Cat

Macaque



The networks of cortical connections in mammalian brain have been shown to have small-world structural properties



Our analysis reveals their dynamical properties to be consistent with modular “small-world” networks

Fast synchronization of neuronal activity within a module :
The mechanism for efficient neural information processing ?

How about other kinds of mesoscopic structures ?

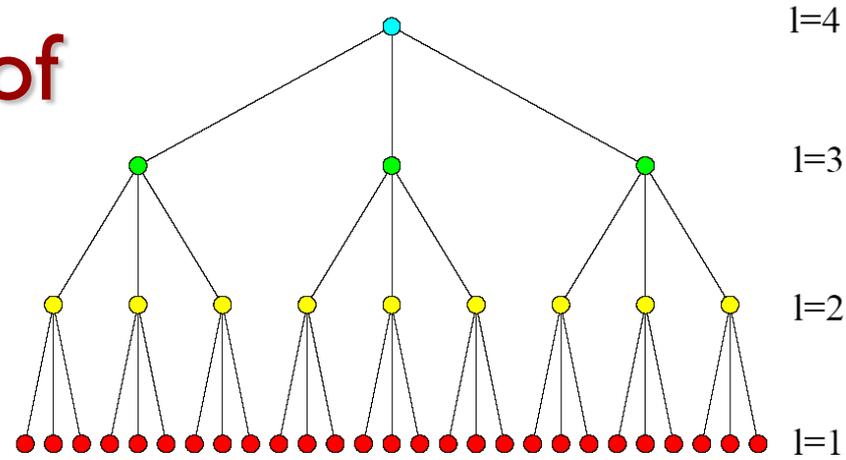
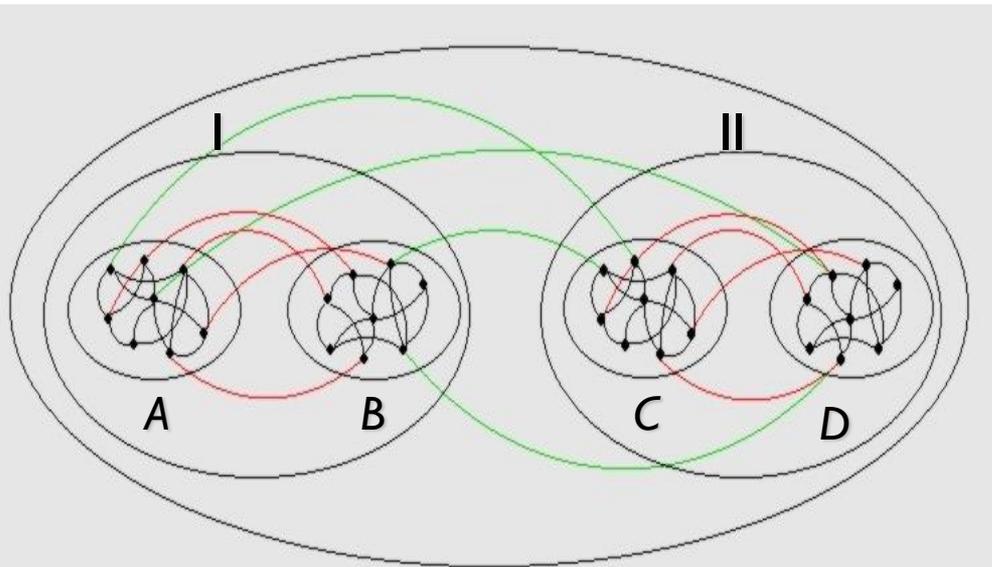
E.g., Hierarchy

Hierarchical Modular networks

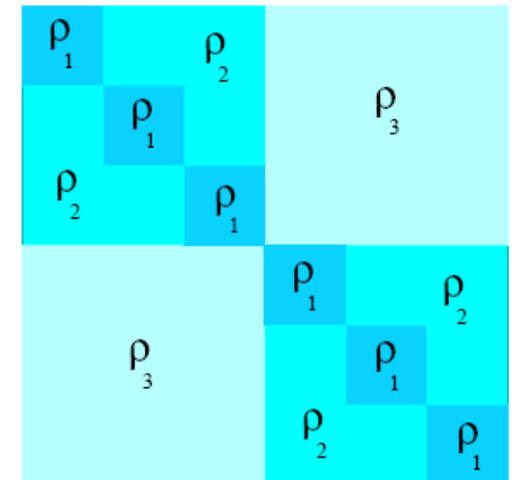
Modules may occur at different levels of hierarchy

Level 1: Modules A, B, C, D

Level 2: Meta-Modules I, II



- $r = 1$: randomly coupled network.
- $r = 0$: isolated sub-networks (modules)
- $0 < r < 1$: hierarchically structured network.

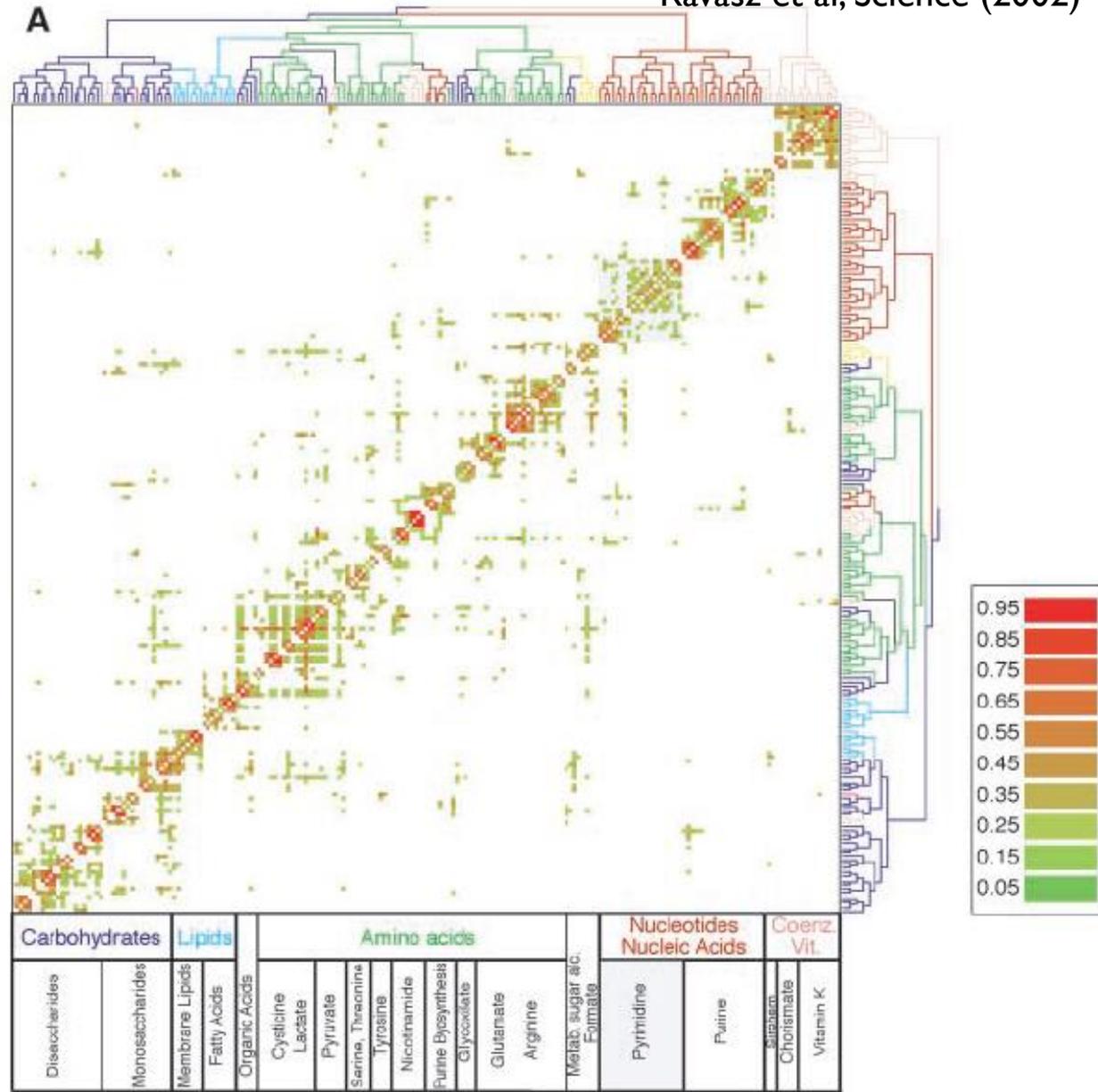
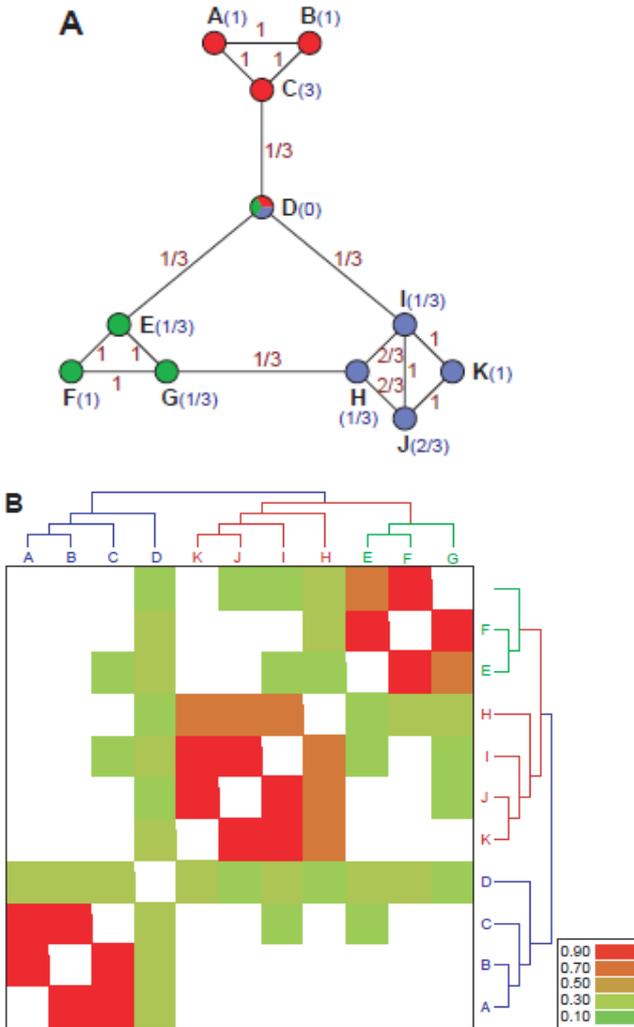


Hierarchical modularity in metabolic network

Topological overlap:

$$O_T(i, j) = \frac{\sum_{l=1}^N l_{i,l} \cdot l_{j,l} + l_{i,j}}{\min(k_i, k_j) + 1 - l_{i,j}}$$

Ravasz et al, Science (2002)



Hierarchical Modular Networks exhibit several distinct time-scales – equal to the number of hierarchical levels (Sinha & Poria, 2011)

Synchronization of phase oscillators in hierarchical modular network show as many distinct time-scales as number of hierarchical levels ...
Reflected in the eigenvalue spectra

